A Decomposition Algorithm for the Sparse Generalized Eigenvalue Problem

Ganzhao Yuan^{1,3}, Li Shen², Wei-Shi Zheng³

¹ Center for Quantum Computing, Peng Cheng Laboratory ² Tencent AI Lab ³ Sun Yat-sen University

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- The Proposed Decomposition Algorithm
- Existing Sparse Optimization Methods
- Theoretical Analysis
- Experiments

The Sparse Generalized Eigenvalue Problem

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The Sparse Generalized Eigenvalue Problem

$$\min_{\mathbf{x}\neq\mathbf{0},\|\mathbf{x}\|_{0}\leq s}f(\mathbf{x})\triangleq\frac{\mathbf{x}^{T}\mathbf{A}\mathbf{x}}{\mathbf{x}^{T}\mathbf{C}\mathbf{x}}.$$

- Statistical learning models
 - Principle Component Analysis (PCA)

$$\min_{x\neq 0} \frac{-x^{T}\Sigma x}{x^{T}x}$$

Pisher Discriminant Analysis (FDA)

$$\mathsf{min}_{\mathsf{x}\neq 0} ~ \tfrac{-\mathsf{x}^{\intercal}((\mu_{(1)}-\mu_{(2)})(\mu_{(1)}-\mu_{(2)})^{\intercal})\mathsf{x}}{\mathsf{x}^{\intercal}(\Sigma_{(1)}+\Sigma_{(2)})\mathsf{x}}$$

Onter Control Correlation Analysis (CCA)

$$\min_{\mathbf{x}} \frac{-\mathbf{x}^{\mathcal{T}} \begin{pmatrix} \mathbf{0} & \boldsymbol{\Sigma}_{\mathbf{x}\mathbf{y}} \\ \boldsymbol{\Sigma}_{\mathbf{y}\mathbf{x}} & \mathbf{0} \end{pmatrix} \mathbf{x}}{\mathbf{x}^{\mathcal{T}} \begin{pmatrix} \boldsymbol{\Sigma}_{\mathbf{x}\mathbf{x}} & \mathbf{0} \\ \mathbf{0} & \boldsymbol{\Sigma}_{\mathbf{y}\mathbf{y}} \end{pmatrix} \mathbf{x}}$$

• Applications: object recognition, object dection, visual tracking, pixel selection, text summarization

The Proposed Algorithm

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The Concept of Block-k Optimality

• Optimization Problem:

$$\min_{\mathbf{x}\in\mathbb{R}^n} F(\mathbf{x})$$

- Variable Splitting: $B \in \mathbb{N}^k$ is a subset of $\{1, 2, ..., n\}$. $N \triangleq \{1, ..., n\} \setminus B$, $\mathbf{x} = \mathbf{I}\mathbf{x} = (\mathbf{U}_B\mathbf{U}_B^T + \mathbf{U}_N\mathbf{U}_N^T)\mathbf{x} = \mathbf{U}_B\mathbf{x}_B + \mathbf{U}_N\mathbf{x}_N$.
- Block-k Optimal Solution

$$\mathcal{P}(B, \mathbf{x}) \triangleq \arg\min_{\mathbf{x}_B} F(\mathbf{U}_B \mathbf{x}_B + \mathbf{U}_N \mathbf{x}_N)$$

 $\{\bar{\mathbf{x}}_B = \mathcal{P}(B, \bar{\mathbf{x}}) \text{ for all } |B| = k\} \Leftrightarrow \bar{\mathbf{x}} \text{ is the block-k optimal solution}$

Block-k Optimality Measure [{B_(i)}^{C^k_n}_{i=1} denotes all the possible combinations]

$$\mathcal{M}(\mathbf{x}) \triangleq \frac{1}{C_n^k} \sum_{i=1}^{C_n^k} \|\mathcal{P}(\mathcal{B}_{(i)}, \mathbf{x}) - \mathbf{x}_{\mathcal{B}_{(i)}}\|_2^2$$

 $\mathcal{M}(\bar{x}) = 0 \Leftrightarrow \bar{x}$ is the block-k optimal solution

Optimization Problem:

$$\min_{\mathbf{x}\neq\mathbf{0}, \|\mathbf{x}\|_{0}\leq s} f(\mathbf{x}) \triangleq \frac{\frac{1}{2}\mathbf{x}^{T}\mathbf{A}\mathbf{x}}{\frac{1}{2}\mathbf{x}^{T}\mathbf{C}\mathbf{x}}$$

We define $h(\mathbf{x}) \triangleq \frac{1}{2}\mathbf{x}^{T}\mathbf{A}\mathbf{x}, \ g(\mathbf{x}) \triangleq \frac{1}{2}\mathbf{x}^{T}\mathbf{C}\mathbf{x}.$

$$\begin{split} h(\mathbf{x}_B, \mathbf{x}_N) &= \frac{1}{2} \mathbf{x}_B^T \mathbf{A}_{BB} \mathbf{x}_B + \frac{1}{2} \mathbf{x}_N^T \mathbf{A}_{NN} \mathbf{x}_N + \langle \mathbf{x}_B, \mathbf{A}_{BN} \mathbf{x}_N \rangle, \\ g(\mathbf{x}_B, \mathbf{x}_N) &= \frac{1}{2} \mathbf{x}_B^T \mathbf{C}_{BB} \mathbf{x}_B + \frac{1}{2} \mathbf{x}_N^T \mathbf{C}_{NN} \mathbf{x}_N + \langle \mathbf{x}_B, \mathbf{C}_{BN} \mathbf{x}_N \rangle. \end{split}$$

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Input: k, θ , a feasible solution x^0 , t = 0while not converge do

S1 Find a working set *B* of size *k*. Denote $N \triangleq \{1, ..., n\} \setminus B$. **S2** Solve the following subproblem *globally*:

$$\mathbf{x}_{B}^{t+1} \leftarrow \arg\min_{\mathbf{x}_{B}} \frac{h(\mathbf{x}_{B}, \mathbf{x}_{N}^{t}) + \frac{\theta}{2} \|\mathbf{x}_{B} - \mathbf{x}_{B}^{t}\|_{2}^{2}}{g(\mathbf{x}_{B}, \mathbf{x}_{N}^{t})}$$

s.t. $\|\mathbf{x}_{B}\|_{0} + \|\mathbf{x}_{N}^{t}\|_{0} \leq s$ (1)

S2 Increment *t* by 1 end

Algorithm 1: The Proposed Decomposition Algorithm

- A new proximal strategy (only applied to the numerator) \Rightarrow sufficient descent and global convergence
- **2** When k = n, the subproblem reduces to the original problem.
- Sinding the working set
 - Random strategy. Select one combination (which contains k coordinates) from the whole working set {B_(i)}^{C_n}_{i=1} uniformly.
 - ② Swapping strategy. Pick the top pairs of coordinates that lead to the greatest descent by measuring D ∈ ℝ^{|S(x)|×|Z(x)|}:

$$\mathbf{D}_{i,j} = \min_{\beta} f(\mathbf{x}^t + \beta \mathbf{e}_i - \mathbf{x}_j^t \mathbf{e}_j) - f(\mathbf{x}^t).$$

with $\mathcal{S}(\mathbf{x}) \triangleq \{i \mid \mathbf{x}_i \neq 0\}$ and $\mathcal{Z}(\mathbf{x}) \triangleq \{j \mid \mathbf{x}_j = 0\}.$

Remarks on the Decomposition Algorithm

The subproblem reduces to the following problem:

$$\min_{\mathbf{z}\in\mathbb{R}^{k}, \|\mathbf{z}\|_{0}\leq q} p(\mathbf{z}) \triangleq \frac{\frac{1}{2}\mathbf{z}^{T}\bar{\mathbf{Q}}\mathbf{z} + \bar{\mathbf{p}}^{T}\mathbf{z} + \bar{w}}{\frac{1}{2}\mathbf{z}^{T}\bar{\mathbf{R}}\mathbf{z} + \bar{\mathbf{c}}^{T}\mathbf{z} + \bar{v}}$$

Our solution:

• We consider the following problem:

 $\min_{\mathbf{z}\in\mathbb{R}^{k}} p(\mathbf{z}), \ s.t. \ \mathbf{z}_{K} = \mathbf{0}$

where K has $\sum_{i=0}^{q} C_k^i$ possible choices. • It reduces to the following problem:

$$\min_{\mathbf{y}} \mathcal{L}(\mathbf{y}) \triangleq \frac{\frac{1}{2}\mathbf{y}^{\mathsf{T}}\mathbf{Q}\mathbf{y} + \mathbf{p}^{\mathsf{T}}\mathbf{y} + w}{\frac{1}{2}\mathbf{y}^{\mathsf{T}}\mathbf{R}\mathbf{y} + \mathbf{c}^{\mathsf{T}}\mathbf{y} + v}$$

Solving the Quadratic Fractional Problem Globally

$$\min_{\mathbf{y}} \mathcal{L}(\mathbf{y}) \triangleq \frac{\frac{1}{2}\mathbf{y}^{T}\mathbf{Q}\mathbf{y} + \mathbf{p}^{T}\mathbf{y} + w}{\frac{1}{2}\mathbf{y}^{T}\mathbf{R}\mathbf{y} + \mathbf{c}^{T}\mathbf{y} + v}$$
$$\mathcal{J}(\alpha) = 0, \text{ with } \mathcal{J}(\alpha) \triangleq \min_{\mathbf{y}} u(\mathbf{y}) - \alpha q(\mathbf{y})$$

We have the following results.

- It holds that: $\lambda_{\min}(\mathsf{Z}) \leq \min_{\mathsf{y}} \mathcal{L}(\mathsf{y}) < \lambda_{\min}(\mathsf{O}).$
- The dual equivalent function $\mathcal{J}(\alpha)$ is monotonically decreasing on the range $\lambda_{\min}(\mathbf{Z}) \leq \alpha < \lambda_{\min}(\mathbf{O})$.
- The optimal solution $\mathbf{y}^* = \mathbf{R}^{-1/2}(\mathbf{u}^* \mathbf{R}^{-1/2}\mathbf{c})$, with $\mathbf{u}^* = -(\mathbf{O} \alpha^*\mathbf{I})^{-1}\mathbf{g}$ and α^* being the unique root of the equation $\mathcal{J}(\alpha) = 0$ on the range $\lambda_{\min}(\mathbf{Z}) \le \alpha < \lambda_{\min}(\mathbf{O})$.

Contributions of this paper:

- A decomposition algorithm, it finds stronger stationary points.
- Iwo strategies to find the working set
- Two methods to solve the subproblem
- A convergence analysis for 'DEC'
- OEC' consistently outperforms existing methods

Existing Sparse Optimization Methods

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Optimization Problem:

$$\min_{\mathbf{x}\in\mathbb{R}^n} f(\mathbf{x}), \ s.t. \ \|\mathbf{x}\|_0 \leq s$$

- f(x): smooth, convex, its gradient is *L*-Lipschitz continuous
- Existing methods
 - Relaxed approximation method
 - ② Greedy pursuit method
 - Ombinatorial search method
 - Gradient projection method

Relaxed approximation method

- convex: $f(\mathbf{x}) + \lambda \|\mathbf{x}\|_1$, $f(\mathbf{x}) + \lambda \|\mathbf{x}\|_{\mathsf{tok}-k}$
- nonconvex: $f(\mathbf{x}) + \lambda \|\mathbf{x}\|_p$, reweighted ℓ_1 norm
- \Rightarrow our method directly controls the sparsity of the solution
- ② Greedy pursuit method
 - the solution MUST be initialized to zero
 - $S = \emptyset$, $S = S \cup i_1$, $\min_{\mathbf{x}_s} f(\mathbf{x})$, $S = S \cup i_2$, $\min_{\mathbf{x}_s} f(\mathbf{x})$, ...

 \Rightarrow our method is a greedy coordinate descent algorithm

Ombinatorial search method: global optimization methods

- cutting plane methods
- branch-and-cut methods

 \Rightarrow our method leverages the effectiveness of combinatorial search method methods

Gradient projection method

•
$$\mathbf{x}^{k+1} = \operatorname{Proj}_{s}(\mathbf{x}^{k} - \gamma \nabla f(\mathbf{x}^{k}))$$

• with
$$\operatorname{Proj}_{s}(\mathbf{a}) = \arg\min_{\mathbf{x}} \frac{1}{2} \|\mathbf{x} - \mathbf{a}\|_{2}^{2}, \ s.t. \ \|\mathbf{x}\|_{0} \leq s$$

 \Rightarrow our method significantly outperforms gradient projection method

Theoretical Analysis

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Basic Stationary Point

 $\breve{\mathbf{x}}$ is called a basic stationary point if the following holds. $\breve{\mathbf{x}} = \arg\min_{\mathbf{y}} f(\mathbf{y}), \ s.t. \ \mathbf{y}_{Z} = \mathbf{0}, \ |S| \le s$, where $Z \triangleq \{i | \breve{\mathbf{x}}_{i} = 0\}$ and $S \triangleq \{i | \breve{\mathbf{x}}_{i} \ne 0\}$.

L-Stationary Point

A solution $\hat{\mathbf{x}}$ is an *L*-stationary point if it holds that: $\hat{\mathbf{x}} = \operatorname{Proj}_{s}(\hat{\mathbf{x}} - \nabla f(\hat{\mathbf{x}})/L).$

Block-k Stationary Point

A solution $\bar{\mathbf{x}}$ is a block-*k* stationary point if it holds that: $\bar{\mathbf{x}} \in \arg\min_{\mathbf{z} \in \mathbb{R}^n} \mathcal{P}(\mathbf{z}; \bar{\mathbf{x}}, B) \triangleq \{F(\mathbf{z}), s.t. \mathbf{z}_N = \bar{\mathbf{x}}_N\}, \forall |B| = k, N \triangleq \{1, ..., n\} \setminus B.$

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Optimality Hierarchy



Relations between the three types of stationary point.

We have the following optimality hierarchy ^a:

^aGanzhao Yuan, Li Shen, Wei-Shi Zheng. A Hybrid Method of Combinatorial Search and Coordinate Descent for Discrete Optimization. arXiv preprint, 2017. url: https://arxiv.org/abs/1706.06493.

Optimization Problems:

$$\min_{\mathbf{x}\in\mathbb{R}^{n}} \ \frac{1}{2}\mathbf{x}^{T}\mathbf{Q}\mathbf{x} + \mathbf{x}^{T}\mathbf{p}, \ s.t. \ \|\mathbf{x}\|_{0} \le s$$
$$n = 6, \ \mathbf{Q} = \mathbf{c}\mathbf{c}^{T} + \mathbf{I}, \ \mathbf{p} = \mathbf{1}, \ \mathbf{c} = [1 \ 2 \ 3 \ 4 \ 5 \ 6]^{T}, \ s = 4.$$

Number of points satisfying optimality conditions.

Basic-	L-Stat.	Block-1	Block-2	Block-3	Block-4	Block-5	Block-6
Stat.		Stat.	Stat.	Stat.	Stat.	Stat.	Stat.
57	56		3	1	1	1	1

Global Convergence Properties.

Assume that the subproblem is solved globally. We have the following results.

- When the random strategy is used to find the working set, we have lim_{t→∞} 𝔼[||x^{t+1} x^t||] = 0 and Algorithm 1 converges to the block-k stationary point in expectation.
- When the swapping strategy is used to find the working set with k ≥ 2, we have lim_{t→∞} $||\mathbf{x}^{t+1} \mathbf{x}^t|| = 0$ and Algorithm 1 converges to the block-2 stationary point deterministically.

Experiments

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Optimization problems:

$$\min_{\mathbf{x}} \frac{\mathbf{x}^{T} \mathbf{A} \mathbf{x}}{\mathbf{x}^{T} \mathbf{C} \mathbf{x}}, \ s.t. \ \|\mathbf{x}\|_{0} \leq s$$

Compared Methods:

- Truncated Power Method (TPM) [Yuan & Zhang, JMLR2013]
- Coordinate-Wise Algorithm (CWA) [Beck & Vaisbourd, JOTA2016]
- Truncated Rayleigh Flow (TRF) [Tan, et al., JRSS2018]
- Quadratic Majorization Method (QMM) [Song, et al., TIP2015]
- Proposed Decomposition Method ¹ (DEC-R*i*-G*j*) [This Paper]

Ganzhao Yuan^{1,3}, Li Shen², Wei-Shi Zheng³

Convergence Behavior: i



Conclusions:

{TPM,CWA,TRF} converge faster, but result in poor accuracy.

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- DEC-B(R10S0) achieves a lower objective value than DEC-B(R6S0). Larger k implies stronger stationary points.
- The swapping strategy plays an indispensable role.

Experimental Results



Conclusions:

- CWA is not stable (much worse results on 'w1a').
- ② DEC achieves lowest objective values.
- South DEC-B and DEC-C perform similarly.

Computational Efficiency



Conclusions:

DEC takes less than 15 seconds to converge in all our instances.

Image: Image:

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② DEC is practical and it is much more efficient than QMM.

Thank You!

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